Docket No. AT9-98-303

METHOD AND APPARATUS TO RETAIN APPLET SECURITY

PRIVILEGES OUTSIDE OF THE JAVA VIRTUAL MACHINE

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BACKGROUND OF THE INVENTION

1. Technical Field:

10   The present invention relates to a system for decoupling a java virtual

machine (JVM) from a browser. Specifically, the invention relates to a system to

retain applet security privileges outside of the JVM, thereby allowing internet

browser decoupling from the JVM.

2. Description of Related Art:

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The Internet, also referred to as an "internetwork", in communications is a set of

computer networks, possibly dissimilar, joined together by means of gateways that

handle data transfer and the conversion of messages from the sending network to the

protocols used by the receiving network (with packets if necessary). When capitalized,

20   the term "Internet" refers to the collection of networks and gateways that use the TCP/IP

suite of protocols.

Docket No. AT9-98-303

The Internet has become a cultural fixture as a source of both information and entertainment. Many businesses are creating Internet sites as an integral part of their marketing efforts, informing consumers of the products or services offered by the business or providing other information seeking to engender brand loyalty. Many

5   federal, state, and local government agencies are also employing Internet sites for informational purposes, particularly agencies which must interact with virtually all segments of society such as the Internal Revenue Service and secretaries of state. Operating costs may be reduced by providing informational guides and/or searchable databases of public records online.

10 Currently, the most commonly employed method of transferring data over the Internet is to employ the World Wide Web environment, also called simply "the web". Other Internet resources exist for transferring information, such as File Transfer Protocol (FTP) and Gopher, but have not achieved the popularity of the web. In the web environment, servers and clients effect data transaction using the

15 Hypertext Transfer Protocol (HTTP), a known protocol for handling the transfer of various data files (e.g., text, still graphic images, audio, motion video, etc.). Information is formatted for presentation to a user by a standard page description language, the Hypertext Markup Language (HTML). In addition to basic presentation formatting, HTML allows developers to specify "links" to other web

20 resources identified by a Uniform Resource Locator (URL). A URL is a special syntax identifier defining a communications path to specific information. Each logical block of information accessible to a client, called a "page" or a "web page", is

Docket No. AT9-98-303

identified by a URL. The URL provides a universal, consistent method for finding

and accessing this information by the web "browser". A browser is a program

capable of submitting a request for information identified by a URL at the client

machine. Retrieval of information on the web is generally accomplished with an

5   HTML-compatible browser.

When a user desires to retrieve a page, a request is submitted to a server

connected to a client computer at which the user is located and may be handled by a

series of servers to effect retrieval of the requested information. The information is

provided to the client formatted according to HTML. Typically, personal computers

10   (PCs) along with work stations are typically used to access the Internet.

Often applications or programs may be sent to a computer from a web server

across the Internet. Java applications are becoming increasingly more prevalent as the

type of application sent between web servers and client computers. Java applications are

common on the Internet and becoming more increasingly common in intranets and in

15   other types of networks used in businesses.

Java is an object oriented programming language and environment focusing

on defining data as objects and the methods that may be applied to those objects. Java

supports only a single inheritance, meaning that each class can inherit from only one

other class at any given time. Java also allows for the creation of totally abstract

20   classes known as interfaces, which allow the defining of methods that may be shared

with several classes without regard for how other classes are handling the methods.

Docket No. AT9-98-303

The Java virtual machine (JVM) is a virtual computer component that resides

only in memory. The JVM allows Java programs to be executed on a different

platform as opposed to only the one platform for which the code was compiled. Java

programs are compiled for the JVM. In this manner, Java is able to support

5 applications for many types of data processing systems, which may contain a variety

of central processing units and operating systems architectures. To enable a Java

application to execute on different types of data processing systems, a compiler

typically generates an architecture-neutral file format – the compiled code is

executable on many processors, given the presence of the Java run-time system. The

10 Java compiler generates bytecode instructions that are non-specific to a particular

computer architecture. A bytecode is a machine independent code generated by the

Java compiler and executed by a Java interpreter. A Java interpreter is a part in the

JVM that alternately decodes and interprets a bytecode or bytecodes. These bytecode

instructions are designed to be easy to interpret on any computer and easily translated

15 on the fly into native machine code.

Many currently available web browsers, such as Netscape Communicator,

which is available from Netscape Communications Corporation, incorporate fixed,

embedded JVMs in which the browsers pass fixed options to the JVM. Presently,

however, JVMs updates are provided more often than web browser updates. Such a

20 situation prevents users from taking advantage of improved versions of JVMs until

the web browser is updated. Therefore, it would be advantageous to have an

improved method and apparatus for providing users an ability to use more recent

Docket No. AT9-98-303

versions of JVMs without having to wait for an updated version of the web browser.

Specifically, a fixed JVM within the browser limits the flexibility of developers and

users of Java applets within the browser. "Applets" are mini applications that

typically run inside a Java-enabled browser.

5   Netscape implements a proprietary Java security scheme within their browser

(and JVM). Applets can request, and subsequently be granted certain privileges, such

as being able to read or write to the local hard drive. These privileges, when granted,

are stored in the JVM's call stack frame which corresponds to the method which has

been granted access. These privileges remain in effect until the method (function)

10   which requested it goes out of scope. In other words, when the function is popped

from the call stack. These privileges are also removed when a user explicitly reverts

the privilege.

          The implementation of storing privileges on the JVM's stack frame requires

JVM modifications for its implementation, since the stack frame data structure is

15   internal to the JVM. This implementation prevents the usage of a de-integrated JVM.

Netscape has also changed many system Java classes to explicitly code enable

Privilege/revert Privilege calls. By doing so, a copy of a Java source file for a given

level of Java is modified with these calls to enable or revert privileges. This locks the

browser into a particular level of Java source.

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Docket No. AT9-98-303

Therefore, a need exists for a method of decoupling the JVM from the browser. Once decoupled, a more up-to-date version of the JVM can be utilized. However, applet security privileges must be retained outside of the JVM by such a system.

Docket No. AT9-98-303

SUMMARY OF THE INVENTION

The present invention addresses the need to decouple the JVM from the browser while retaining applet security privileges. Each Java thread that enables a privilege will now create an entry that describes the privilege in a linked list based on 5 the stack frame address. Sufficient information is stored in the link list entry for validation purposes, to guard against the case where a function that was granted privilege has since returned and the same stack frame is being reused by another function which should not have privilege. The information that is kept is the method name and signature (parameter list) of the function requesting the privilege, as well as 10 the calling function's return address (instruction pointer). When a query is done to determine if the applet has the required privilege, the linked list for its thread will be searched and the privilege (if found and valid) will be returned. Privileges are reverted by removing the element from the linked list.

System Java classes also no longer need to be modified with enable/revert 15 privilege calls. Instead, system classes that require specific privileges will have them implicitly granted. An implicit privilege list is created in memory. The implicit privilege list can map system classes, functions within them that require privileges, and the privileges that are needed. When a request is made to determine if a function has the required privilege, the linked list of privileges will first be checked. If no 20 privilege is found, a subsequent search of the implicit privilege list will determine whether the requesting function is allowed access to the resource. This apparatus can

Docket No. AT9-98-303

also be used in future implementations for storing additional Java stack frame-based

information.

BRIEF DESCRIPTION OF THE DRAWINGS

The novel features believed characteristic of the invention are set forth in the appended claims. The invention itself, however, as well as a preferred mode of use, further objectives and advantages thereof, will best be understood by reference to the

5   following detailed description of an illustrative embodiment when read in conjunction with the accompanying drawings, wherein:

Figure 1 is a pictorial representation of a distributed data processing system in which the present invention may be implemented;

Figure 2 is a block diagram of a data processing system that may be

10   implemented as a server;

Figure 3 is a block diagram of a data processing system in which the present invention may be implemented;

Figure 4 is a block diagram illustrating the stack frame shadow apparatus used in the implementation of the present invention;

15   Figure 5 is a block diagram illustrating the Java stack frame data structure and the Java stack frame extension data structure;

Figure 6 is a flow chart of the method of setting a Java stack frame extension in the apparatus of Figure 4;

Figure 7 is a flow chart of the method of querying to get a Java stack frame

20   extension in the apparatus; and

Figure 8 is a flow chart of the method for removing a Java stack frame extension in the apparatus.

Docket No. AT9-98-303

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

With reference now to the figures, and in particular with reference to **Figure 1**, a pictorial representation of a distributed data processing system in which the present invention may be implemented is depicted.

5   Distributed data processing system **100** is a network of computers in which the present invention may be implemented. Distributed data processing system **100** contains a network **102**, which is the medium used to provide communications links between various devices and computers connected together within distributed data processing system **100**. Network **102** may include permanent connections, such as wire

10   or fiber optic cables, or temporary connections made through telephone connections.

In the depicted example, a server **104** is connected to network **102** along with storage unit **106**. In addition, clients **108, 110**, and **112** also are connected to a network **102**. These clients **108, 110**, and **112** may be, for example, personal computers of network computers. For purposes of this application, a network computer is any

15   computer, coupled to a network, which receives a program or other application from another computer coupled to the network. In the depicted example, server **104** provides data, such as boot files, operating system images, and applications to NCs **108-112**. NCs **108, 110**, and **112** are clients to server **104**. Distributed data processing system **100** may include additional servers, clients, and other devices not shown.

20   In the depicted example, distributed data processing system **100** is the Internet with network **102** representing a worldwide collection of networks and gateways that use the TCP/IP suite of protocols to communicate with one another. At the heart of the

Docket No. AT9-98-303

Internet is a backbone of high-speed data communication lines between major nodes or host computers, consisting of thousands of commercial, government, educational, and other computer systems, that route data and messages. Of course, distributed data processing system **100** also may be implemented as an umber of different types of

5   networks, such as for example, an intranet or a local area network.

Figure 1 is intended as an example, and not as an architectural limitation for the processes of the present invention.

Referring to **Figure 2**, a block diagram of a data processing system which may be implemented as a server, such as server **104** in **Figure 1**, is depicted in accordance to

10   the present invention. Data processing system **200** may be a symmetric multiprocessor (SMP) system including a plurality of processors **202** and **204** connected to system bus **206**. Alternatively, a single processor system may be employed. Also connected to system bus **206** is memory controller/cache **208**, which provides an interface to local memory **209**. I/O bus bridge **210** is connected to system bus **206** and provides an

15   interface to I/O bus **212**. Memory controller/cache **208** and I/O bus bridge **210** may be integrated as depicted.

Peripheral component interconnect (PCI) bus bridge **214** connected to I/O bus **212** provides an interface to PCI local bus **216**. A number of modems **218-220** may be connected to PCI bus **216**. Typical PCI bus implementations will support four PCI

20   expansion slots or add-in connectors. Communications links to network computers **108-112** in **Figure 1** may be provided through modem **218** and network adapter **220** connected to PCI local bus **216** through add-in boards.

Additional PCI bus bridges **222** and **224** provide interfaces for additional PCI

buses **226** and **228**, from which additional modems or network adapters may be

supported. In this manner, server **200** allows connections to multiple network

computers. A memory mapped graphics adapter **230** and hard disk **232** may also be

5   connected to I/O bus **212** as depicted, either directly or indirectly.

Those of ordinary skill in the art will appreciate that the hardware depicted in

Figure 2 may vary. For example, other peripheral devices, such as optical disk drive

and the like also may be used in addition or in place of the hardware depicted. The

depicted example is not meant to imply architectural limitations with respect to the

10   present invention.

The data processing system depicted in **Figure 2** may be, for example, an IBM

RISC/System 6000 system, a product of International Business Machines Corporation in

Armonk, New York, running the Advanced Interactive Executive (AIX) operating

system.

15   With reference now to **Figure 3**, a block diagram of a data processing system

300 in which the present invention may be implemented is illustrated. Data processing

system **300** is an example of a client computer. Data processing system **300** employs a

peripheral component interconnect (PCI) local bus architecture. Although the depicted

example employs a PCI bus, other bus architectures such as Micro Channel and ISA

20   may be used. Processor **302** and main memory **304** are connected to PCI local bus **306**

through PCI bridge **308**. PCI bridge **308** also may include an integrated memory

controller and cache memory for processor **302**. Additional connections to PCI local

Docket No. AT9-98-303

bus **306** may be made through direct component interconnection or through add-in

boards. In the depicted example, local area network (LAN) adapter **310**, SCSI host bus

adapter **312**, and expansion bus interface **314** are connected to PCI local bus **306** by

direct component connection. In contrast, audio adapter **316**, graphics adapter **318**, and

5   audio/video adapter (A/V) **319** are connected to PCI local bus **306** by add-in boards

inserted into expansion slots. Expansion bus interface **314** provides a connection for a

keyboard and mouse adapter **320**, modem **322**, and additional memory **324**. SCSI host

bus adapter **112** provides a connection for hard disk drive **326**, tape drive **328**, and CD-

ROM **330** in the depicted example. Typical PCI local bus implementations will support

10   three or four PCI expansion slots or add-in connectors.

An operating system runs on processor **302** and is used to coordinate and

provide control of various components within data processing system **300** in **Figure 1**.

The operating system may be a commercially available operating system such as OS/2,

which is available from International Business Machines Corporation. "OS/2" is a

15   trademark of from International Business Machines Corporation. An object oriented

programming system such as Java may run in conjunction with the operating system and

provides calls to the operating system from Java programs or applications executing on

data processing system **300**. Instructions for the operating system, the object-oriented

operating system, and applications or programs are located on storage devices, such as

20   hard disk drive **326** and may be loaded into main memory **304** for execution by

processor **302.**

Docket No. AT9-98-303

Those of ordinary skill in the art will appreciate that the hardware in **Figure 3**

may vary depending on the implementation. For example, other peripheral devices,

such as optical disk drives and the like may be used in addition to or in place of the

hardware depicted in **Figure 3**. The depicted example is not meant to imply architectural

5   limitations with respect to the present invention. For example, the processes of the

present invention may be applied to multiprocessor data processing system.

Figure 4 is a block diagram illustrating the relationship between the JVM and

the browser and also illustrating the basic steps for creating the linked list. The

system **400** includes a browser **402**. The browser can incorporate or interact with an

10   existing JVM **404**. The existing JVM includes a stack frame for a first thread **406** and

a stack frame for a second thread **408**. As mentioned above, privileges are stored in

the stack frames. Each stack can contain a plurality of frames such as those

designated by the x and y arrows. Under normal operation, the internet browser **402**

queries the existing JVM **404** to get a stack frame (step **410**). In response, the

15   existing JVM returns the appropriate stack frame (step **412**). The present invention

involves the creation of a stack frame shadow apparatus 420. The creation of the

stack frame shadow apparatus involves setting the stack frame extension using the

thread identifier and frame (step **414**). Then, when the browser needs to know if a

privilege is available, the stack frame shadow apparatus 420 is queried to get the

20   frame extension by thread identifier and frame (step **416**). The frame extension is

returned (step **418**) and then the frame extension can be removed from the stack

frame shadow apparatus (step **420**). Several of these steps will be discussed below in

Docket No. AT9-98-303

greater detail. Within the stack frame shadow apparatus 420 [426] are stack frame

extensions for first thread 422 and for the second thread 424.

Figure 5 provides a comparison between the Java stack frame data structure

502 and the Java stack frame extension data structure 510 in the stack frame shadow

apparatus 420 [426] of Figure 4. The Java stack frame data structure 502 includes the

pointer to the last stack frame 504, the local variables in this stack frame 506, and

other information 508. The Java stack frame extension data structure 510 can include

the Java security privilege 512, other variables 514, the method name and signature

for validation 516, and the return address for validation 518. The pointer to the Java

stack frame 520 is used to map this extension entry with the JVM's stack frame 508.

Figure 6 is a flow chart of the method 600 of setting Java stack frame

extension in the apparatus. First, the system must use the thread identifier to get a list

of frame extension entries (step 602). Next, the system searches for the list of frames

for matching frame entry(ies) (step 604). Next, the system decides whether a

matching frame extension entry has been found (step 606). If yes, then the system

will update the entry with the frame extension information and validation information

(step 608). If not, then the system will create a blank entry and fill the blank entry in

with frame extension information and validation information (step 610).

Figure 7 is a flow chart of the method 700 of getting a Java stack frame

extension located in the shadow apparatus 420 [426]. The system first uses the thread

identifier to get a list of frame extension entries (step 702). Next, the system searches

the list of frame extensions for a matching frame entry (step 704). Next, the system

Docket No. AT9-98-303

decides whether a matching frame extension is found (step **706**). If so, the system

must validate the frame with frame extension information (step **708**). If not, the

system returns with a blank or null frame extension information (step **710**). Once

step **708** is completed, the system further decides whether the frame is valid for the

5   particular entry (step **712**). If not, the system returns to step **710**. If however, the

frame is valid in step **712**, then the system returns with the frame extension

information (step **714**).

Figure 8 illustrates the method **800** of removing the Java stack frame

extension in the shadow apparatus *426*. First, the system uses a thread identifier to

10   get a list of frame extension entries (step **802**). Next, the system searches the list of

frame extensions for a matching frame entry (step **804**). If a matching frame

extension is found (step **806**), then the entry is removed from the list (step **808**). If

no match is found, then the system returns (step **810**) to the caller to revert privilege.

It is important to note that while the present invention has been described in

15   the context of a fully functioning data processing system, those of ordinary skill in the

art will appreciate that the processes of the present invention are capable of being

distributed in a form of a computer readable medium of instructions and a variety of

forms and that the present invention applies equally regardless of the particular type

of signal bearing media actually used to carry out the distribution. Examples of

20   computer readable media include recordable-type media such a floppy disc, a hard

disk drive, a RAM, and CD-ROMs and transmission-type media such as digital and

analog communications links.

Docket No. AT9-98-303

The description of the present invention has been presented for purposes of illustration and description, but is not limited to be exhaustive or limited to the invention in the form disclosed. Many modifications and variations will be apparent to those of ordinary skill in the art. The embodiment was chosen and described in order to best explain the principles of the invention the practical application and to enable others of ordinary skill in the art to understand the invention for various embodiments with various modifications as are suited to the particular use contemplated.